

optimization problem

minimize $f_0(x)$
 s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$

feasible $x \in \text{dom } f_0$
infeasible $P = \emptyset$

optimal value
 $P^* = \inf \{f_0(x) \mid f_0(x) \leq 0, h_i(x) = 0\}$
 P^* is optimal

locally optimal
 $f_0(x) = \inf \{f_0(\beta) \mid \beta \in D, \|\beta - x\|_2 \leq R\}$
 $D = \{x \mid \text{dom } f_0 \cap \bigcap_{i=1}^m \{f_i \leq 0, h_i = 0\}\}$

feasibility problem
 find x
 s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$

convex optimization
 minimize $f_0(x)$
 s.t. $f_i(x) \leq 0$ (convex)
 $A^T x = b$ (affine)

quasiconvex optimization
 minimize $f_0(x)$ w.r.t. quasiconvex
 s.t. $f_i(x) \leq 0$
 $A^T x = b$

local & global optima
 minimize $f_0(x)$
 s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$
 $\|x - x^*\|_2 \leq R$

x^* is **locally optimal** when
 $f_0(x) = \inf \{f_0(\beta) \mid \beta \text{ feasible}, \|\beta - x^*\|_2 \leq R\}$

in convex if x^* is local optimal then x^* is global optimal

proof:
 Δ if $\|y - x^*\|_2 > R$ as if $\|y - x^*\|_2 < R$

Δ we let $\theta = \frac{R}{\|y - x^*\|_2} < \frac{1}{2}$

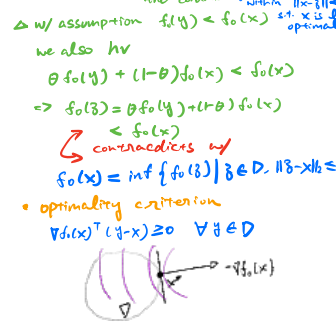
Δ we let $\tilde{x} = (1-\theta)x^* + \theta y$
 convex combination, in convex function

$f_0(\tilde{x}) \leq (1-\theta)f_0(x^*) + \theta f_0(y)$
 $0 < \theta < \frac{1}{2}$

Δ w/ assumption $f_0(y) < f_0(x^*)$ s.t. x^* is **not** optimal

we also hv
 $\theta f_0(y) + (1-\theta)f_0(x^*) < f_0(x^*)$
 $\Rightarrow f_0(\tilde{x}) < f_0(x^*)$
 contradiction w/ $f_0(x^*) = \inf \{f_0(\beta) \mid \beta \in D, \|\beta - x^*\|_2 \leq R\}$

optimality criterion
 $\forall \lambda \in D, \nabla f_0(x^*)^T (y - x^*) \geq 0$



convex problems

Linear programming
 minimize $c^T x + d$
 s.t. $Ax \leq h$
 $Ax = b$

piecewise-linear minimization
 $f_0(x) = \max_{i=1, \dots, m} (a_i^T x + b_i)$

minimize $f_0(x)$
 \Leftrightarrow
 minimize t
 s.t. $A^T x + b_i \leq t, i=1, \dots, m$

Quadratic Program
 minimize $\frac{1}{2} x^T P x + g^T x + r$
 s.t. $Ax \leq h$
 $Ax = b$

$P \in S^n$

Quadratically Constrained Quadratic Program
 minimize $\frac{1}{2} x^T P x + g^T x + r$
 s.t. $\frac{1}{2} x^T P_i x + g_i^T x + r_i \leq 0$
 $Ax = b$
 $P_i \in S^n$ (convex feasible region, what's intersection of ellipsoids + affine set)

Geometric Programming

monomial function
 $f: \mathbb{R}^n \rightarrow \mathbb{R}$ defn: $f(x) = c \cdot x_1^{a_1} \cdot x_2^{a_2} \cdot \dots \cdot x_n^{a_n}$
 $c > 0, a_i \in \mathbb{R}$

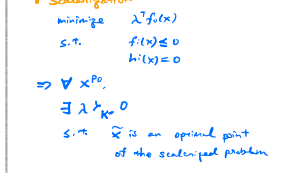
polynomial function
 $f(x) = \sum_{k=1}^K c_k \cdot x_1^{a_{k1}} \cdot x_2^{a_{k2}} \cdot \dots \cdot x_n^{a_{kn}}$
 $c_k > 0, a_{ki} \in \mathbb{R}$

sum of monomial function
 transfer to convex problem
 $y_i = \log x_i, x_i = e^{y_i}$
 $f(x) = f(e^{y_1}, \dots, e^{y_n}) = C(e^{y_1}, \dots, e^{y_n}) = e^{a^T y + b}$ (linear)
 monomial

multicriterion optimization
 minimize $f_0(x) = [f_1(x), \dots, f_m(x)]$
 s.t. $f_i(x) \leq 0$
 $Ax = b$

Pareto optimal
 $(f_0(x) - K) \cap f_0(x) = \emptyset$

Scalarization
 minimize $\lambda^T f_0(x)$
 s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$
 $\Rightarrow \forall x \in P^*, \exists \lambda \succeq_{K^*} 0$
 s.t. \tilde{x} is an optimal point of the scalarized problem



KKT optimality conditions
 let x^* (λ^*, ν^*) be primal dual optimal pnts
 - D - dual gap
 $\therefore x^* = \arg \min_x L(x, \lambda^*, \nu^*)$

$f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^l \nu_i^* h_i(x^*) = 0$

$f_i(x^*) \leq 0$
 $h_i(x^*) = 0$
 $\lambda_i^* \geq 0$
 $\lambda_i^* f_i(x^*) = 0$

\rightarrow KKT optimality conditions!

Second-order Cone Programming

minimize $f^T x$
 $\|Ax + b\|_2 \leq c^T x + d$
 $Fx = g$

SOCP could be understood:
 $\{Ax + b, c^T x + d\} \in \mathbb{R}^k$ lies in SOC

$(Ax + b, c^T x + d) \succeq_{\mathbb{R}^k} 0$
 \Leftrightarrow for cone \mathcal{Q}

e.g. Robust Linear Programming

Duality

The Lagrangian
 minimize $f_0(x)$
 s.t. $f_i(x) \leq 0, i=1, \dots, m$
 $h_i(x) = 0, i=1, \dots, l$

$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^l \nu_i h_i(x)$

The Lagrange Dual Function
 $g(\lambda, \nu) = \inf_{x \in D} L(x, \lambda, \nu)$

$\therefore g(\lambda, \nu) \leq P^* = f_0(x^*)$ (dual feasible)

The Lagrange Dual Problem
 maximize $g(\lambda, \nu)$
 s.t. $\lambda \succeq 0$

Weak Duality
 $D^* \leq P^*$
 best dual value, optimal value, duality gap

Strong Duality
 $D^* = P^*$

Slater's condition
 $\exists x \in \text{relative int } D$
 s.t. $f_i(x) < 0, Ax = b$

optimizing conditions - complementary slackness
 assume strong duality holds:
 $D^* = P^*$
 $g(\lambda^*, \nu^*) = f_0(x^*)$
 $\Rightarrow f_0(x^*) = g(\lambda^*, \nu^*)$
 $= \inf_x (f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^l \nu_i^* h_i(x))$
 $\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^l \nu_i^* h_i(x^*)$
 $= f_0(x^*)$

$\left\{ \begin{array}{l} \text{if } \lambda_i^* > 0 \Rightarrow f_i(x^*) = 0 \\ \text{if } f_i(x^*) < 0 \Rightarrow \lambda_i^* = 0 \end{array} \right.$
 complementary slackness!

$\therefore \sum_{i=1}^m \lambda_i^* f_i(x^*) = 0$

KKT optimality conditions
 let x^* (λ^*, ν^*) be primal dual optimal pnts
 - D - dual gap
 $\therefore x^* = \arg \min_x L(x, \lambda^*, \nu^*)$

$f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^l \nu_i^* h_i(x^*) = 0$

$f_i(x^*) \leq 0$
 $h_i(x^*) = 0$
 $\lambda_i^* \geq 0$
 $\lambda_i^* f_i(x^*) = 0$

\rightarrow KKT optimality conditions!

Perturbation & Sensitivity Analysis

recall
 minimize $f_0(x)$
 s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$

recall
 minimize $f_0(x)$
 s.t. $f_i(x) \leq u_i$
 $h_i(x) = v_i$

recall
 minimize $f_0(x)$
 s.t. $f_i(x) \leq u_i$
 $h_i(x) = v_i$

optimal value (becomes a function)
 w.r.t. (u, v)
 $P^*(u, v) = \inf \{f_0(x) \mid f_i(x) \leq u_i, h_i(x) = v_i\}$

let $(\lambda^*, \nu^*) \rightarrow$ dual optimal pt. of unperturbed problem
 $\Rightarrow P^*(0,0) = g(\lambda^*, \nu^*)$
 $\leq f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^l \nu_i^* h_i(x)$
 $\leq f_0(x) + \lambda^{*T} u + \nu^{*T} v$
 $\Rightarrow P^*(0,0) \leq f_0(x) + \lambda^{*T} u + \nu^{*T} v$
 $\Rightarrow f_0(x) \geq P^*(0,0) - \lambda^{*T} u - \nu^{*T} v$
 $\Rightarrow P^*(u, v) \geq P^*(0,0) - \lambda^{*T} u - \nu^{*T} v$
 $(P^*(u, v) \in \{f_0(x) \mid x \in \text{dom } f_0\})$
 $P^*(u, v) \geq P^*(0,0) - \lambda^{*T} u - \nu^{*T} v$
 perturbation analysis

Sufficiency & Necessity

if P then \mathcal{Q}

$\Delta \mathcal{Q}$ is necessary for P

\mathcal{Q} truth of \mathcal{Q} is guaranteed by truth of P

\mathcal{Q} (satisfying of \mathcal{Q} guarantees satisfying of P)

ΔP is sufficient for P

\mathcal{P} truth of P guarantees truth of \mathcal{Q}

but falsity of P might still hv truth of \mathcal{Q}

如果“要选谁”那“谁合适”
 ① “一定是合适”如果谁谁谁
 ② 如果“未合适”那“一定不合适”
 ③ 如果“要选谁”那“一定合适”
 ④ 如果“不选谁”那“可能合适”

sum:
 if P then \mathcal{Q}
 $\Delta \mathcal{Q}$ is P 's necessary condition
 ΔP is \mathcal{Q} 's sufficient condition

local sensitivity
 assume $P^*(u, v)$ differentiable @ $u=0, v=0$

strong duality holds
 $\Rightarrow P^*(u, v) @ u=0, v=0$
 $\frac{\partial P^*(0,0)}{\partial u_i} = -\lambda_i^*$
 $\frac{\partial P^*(0,0)}{\partial v_i} = -\nu_i^*$

w/ complementary slackness:
 $\lambda_i^* f_i(x^*) = 0$

$f_i(x) < 0, \lambda_i = 0$
 then: value @ $P^*(0,0)$ doesn't change w.r.t. u

$f_i(x) = 0, \lambda_i^* > 0$
 then: value @ $P^*(0,0)$ changes w.r.t. u
 $\lambda_i^* \uparrow \rightarrow$ perturb \uparrow
 $\lambda_i^* \downarrow \rightarrow$ perturb \downarrow

\rightarrow tells us HOW ACTIVE IT IS

Numerical Method of Differential Equations

ODE 4 types

1. separable equations

$\frac{dy}{dx} = P(x)Q(x)$
 e.g. $y' + 2xy = x$
 $\Rightarrow \frac{dy}{dx} + 2xy = x$
 $\Rightarrow \frac{dy}{dx} = x - 2xy$
 $\Rightarrow \frac{dy}{dx} = x(1-2y)$
 $\Rightarrow dy = x(1-2y) dx$
 $\Rightarrow \frac{1}{1-2y} dy = x dx$
 $\Rightarrow \int \frac{1}{1-2y} dy = \int x dx$
 $\Rightarrow -\frac{1}{2} \ln(1-2y) = \frac{1}{2} x^2 + C$
 $\Rightarrow e^{-\ln(1-2y)} = e^{-x^2 + C}$
 $\Rightarrow 1-2y = ce^{-x^2}$
 $\Rightarrow y = \frac{1 - ce^{-x^2}}{2}$

2. homogenous method

$f(kx, ky) = f(x, y)$
 e.g. $\frac{dy}{dx} = \frac{x^2 + y^2}{xy}$
 check $\frac{k^2x^2 + k^2y^2}{kx \cdot ky} = \frac{x^2 + y^2}{xy}$
 let $v = \frac{y}{x} \Rightarrow \frac{1}{x} \frac{dv}{dx} = \frac{v^2 + 1}{v}$
 $\Rightarrow x \frac{dv}{dx} + v = \frac{1+v^2}{v}$
 $\Rightarrow x \frac{dv}{dx} + v = \frac{1}{v} + v$
 $\Rightarrow x \frac{dv}{dx} = \frac{1}{v}$
 $\Rightarrow v dv = \frac{1}{x} dx$
 $\Rightarrow \int v dv = \int \frac{1}{x} dx$
 $\Rightarrow \frac{1}{2} v^2 = \ln(x) + C$
 $\Rightarrow v^2 = 2 \ln(x) + C$
 $v = \pm \sqrt{2 \ln(x) + C}$
 $\frac{y}{x} = \pm \sqrt{2 \ln(x) + C}$
 $y = \pm x \sqrt{2 \ln(x) + C}$

3. Integrating factor

$\frac{dy}{dx} + P(x)y = Q(x)$
 e.g. $\frac{dy}{dx} + 1y = x$
 $P(x) = 1, Q(x) = x$
 $\Rightarrow \frac{dy}{dx} + P(x)y = Q(x)$
 $\mu(x) = e^{\int P(x) dx}$
 $\mu(x) = e^{\int 1 dx} = e^x$
 $\Rightarrow \textcircled{1} \& \textcircled{2}$
 $\mu(x) \left[\frac{dy}{dx} + P(x)y = Q(x) \right]$
 $\Rightarrow e^x \left[\frac{dy}{dx} + y = x \right]$
 $\Rightarrow e^x \frac{dy}{dx} + ye^x = xe^x$
 $\Rightarrow \frac{d}{dx} (e^x y) = xe^x$
 $\Rightarrow \int \frac{d}{dx} (e^x y) = \int xe^x dx$
 $\Rightarrow e^x y = xe^x - e^x + C$
 $\Rightarrow y = x - 1 + \frac{C}{e^x}$
 $= x - 1 + ce^{-x}$

$\frac{dy}{dt} = y - \frac{y^2}{m}$

$\frac{dy}{dt} + P(t)y = Q(t)$
 $\Rightarrow \frac{dy}{dt} + \left(\frac{c}{m}\right)y = y$
 $\Rightarrow \mu(t) = e^{\int \frac{c}{m} dt} = e^{\frac{ct}{m}}$
 $\Rightarrow e^{\frac{ct}{m}} \left(\frac{dy}{dt} + \left(\frac{c}{m}\right)y = y \right)$
 $\Rightarrow e^{\frac{ct}{m}} \frac{dy}{dt} + \frac{c}{m} ye^{\frac{ct}{m}} = ye^{\frac{ct}{m}}$
 $\Rightarrow \frac{d}{dt} (e^{\frac{ct}{m}} y) = ye^{\frac{ct}{m}}$
 $\Rightarrow e^{\frac{ct}{m}} y = \int ye^{\frac{ct}{m}} dt + C$
 $\Rightarrow y = e^{-\frac{ct}{m}} \left[\int ye^{\frac{ct}{m}} dt + C \right]$
 Assume $V(0) = 0$
 $V(0) = \frac{1}{m} y + C' = 0$
 $C' = -\frac{1}{m} y$
 $\therefore v(t) = e^{-\frac{ct}{m}} \left[\int \frac{m}{c} ye^{-\frac{ct}{m}} dt - \frac{1}{m} y \right]$
 $= e^{-\frac{ct}{m}} \frac{m}{c} \left[\int ye^{-\frac{ct}{m}} dt - y \right]$
 $= \frac{m}{c} \left[1 - e^{-\frac{ct}{m}} \right]$

- Linear ODE**
 - Laplace
 - analytically solved
 - $a(x)y + a'(x)y' + \dots + a_n(x)y^{(n)} = b(x)$
 - $1y = f$

Runge-Kutta Methods

Taylor series
 $\frac{dy}{dx} = f(x, y)$
 $\frac{d^{(n)}y}{dx^{(n)}} = f^{(n)}(x, y)$
 $y_{i+1} = y_i + f(x_i, y_i)h$
Euler's Method
 $y_{i+1} = y_i + f(x_i, y_i)h$
 $+ \frac{f'(x_i, y_i)}{1!} h^2 + \frac{f''(x_i, y_i)}{2!} h^3 + \dots + \frac{f^{(n)}(x_i, y_i)}{n!} h^{n+1}$
 $\approx y_i + f(x_i, y_i)h + \frac{f'(x_i, y_i)}{1!} h^2 + \frac{f''(x_i, y_i)}{2!} h^3 + \dots + \frac{f^{(n)}(x_i, y_i)}{n!} h^{n+1}$
 $= y_i + f(x_i, y_i)h$
Runge-Kutta:
 acquire $f(x_i, y_i, h)$
 w/o calculating higher-order term
 $y_{i+1} = y_i + f(x_i, y_i, h)$
 $\Rightarrow f(x_i, y_i, h)$
 $= A_1 k_1 + A_2 k_2 + \dots + A_n k_n + A_{n+1} k_{n+1} + A_{n+2} k_{n+2}$
 $(k_1 = f(x_i, y_i))$
 $(k_2 = f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}f(x_i, y_i)h))$
 $(k_3 = f(x_i + \frac{1}{2}h, y_i + f(x_i, y_i)h))$
 $(k_4 = f(x_i + h, y_i + f(x_i, y_i)h))$
 $A_1, A_2, \dots, A_n, A_{n+1}, A_{n+2} \Rightarrow$ constants
 $n =$ order \Rightarrow **RKN method!**
 e.g.
 $\Delta y = -0.5x^2 + 4x^3 - 10x^4 + 8.5x + 1$
 $\frac{dy}{dx} = -2x^2 + 12x^3 - 20x^4 + 8.5x + 8.5$
 $f(0) = 1$
 $f(0.5) = ?$
RK4:
 $y_{i+1} = y_i + f(x_i, y_i, h)$
 $f(x_i, y_i, h)$
 $= \frac{1}{6} (k_1 + 2k_2 + 2k_3 + k_4)$
 $k_1 = f(x_i, y_i)$
 $k_2 = f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_1 h)$
 $k_3 = f(x_i + \frac{1}{2}h, y_i + k_2 h)$
 $k_4 = f(x_i + h, y_i + k_3 h)$
 \therefore
 $k_1 = f(0) = 8.5$
 $k_2 = f(0 + 0.25, 1 + 0.5 \cdot 8.5) = 4.21875$
 $k_3 = f(0 + 0.25, 1 + 0.5 \cdot 4.21875) = 2.0625$
 $k_4 = f(0.5, 1 + 0.21875) = 1.25$
 $y(0.5) = 1 + \frac{0.5}{6} (8.5 + 2 \cdot 4.21875 + 2 \cdot 2.0625 + 1.25)$
 $= 3.21875 + O(h^5)$

system of equations

$\begin{cases} \frac{dy_1}{dx} = f_1(x, y_1, y_2, \dots, y_n) \\ \frac{dy_2}{dx} = f_2(x, y_1, y_2, \dots, y_n) \\ \vdots \\ \frac{dy_n}{dx} = f_n(x, y_1, y_2, \dots, y_n) \end{cases}$
 e.g.
 $\Delta \begin{cases} \frac{dy_1}{dx} = -0.5y_1 \\ \frac{dy_2}{dx} = 4 - 0.3y_2 - 0.1y_1 \\ y_1 = 4 \\ y_2 = 6 \end{cases} @ x=0$

$\Delta y' = f(x)$
 $= \begin{bmatrix} \frac{dy_1}{dx} \\ \frac{dy_2}{dx} \\ \frac{dy_3}{dx} \end{bmatrix} = \begin{bmatrix} -0.5 & 0 \\ 0 & -0.3 \\ 0 & -0.1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} + \begin{bmatrix} 0 \\ 4 \\ 6 \end{bmatrix}$

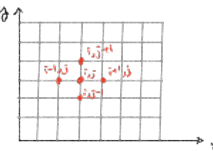
Δ RK4

$y_{i+1} = y_i + \frac{1}{6} (k_1 + 2k_2 + 2k_3 + k_4)h$
 $k_1 = f(x_i, y_i)$
 $k_2 = f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_1 h)$
 $k_3 = f(x_i + \frac{1}{2}h, y_i + k_2 h)$
 $k_4 = f(x_i + h, y_i + k_3 h)$
 $\Delta k_{1,1} = f_1(0, 4) = -0.5 \cdot 4 = -2$
 $k_{1,2} = f_2(0, 6) = 4 - 0.3 \cdot 6 - 0.1 \cdot 4 = 1.8$
 $\Delta k_{2,1} = f_1(0 + \frac{1}{2} \cdot 0.5, 4 + \frac{1}{2} \cdot (-2) \cdot 0.5)$
 $= -0.5 \cdot 3.5 = -1.75$
 $k_{2,2} = f_2(0 + \frac{1}{2} \cdot 0.5, 6 + \frac{1}{2} \cdot 1.8 \cdot 0.5)$
 $= 4 - 0.3 \cdot 6.45 - 0.1 \cdot 3.5 = 1.715$
 $\Delta k_{3,1} = f_1(0 + \frac{1}{2} \cdot 0.5, 4 + 1 \cdot (-1.75) \cdot 0.5)$
 $= -0.5 \cdot 3.5625 = -1.78125$
 $k_{3,2} = f_2(0 + \frac{1}{2} \cdot 0.5, 6 + 1 \cdot 1.715 \cdot 0.5)$
 $= 4 - 0.3 \cdot 6.8575 - 0.1 \cdot 3.5125 = 1.715125$
 $\Delta k_{4,1} = f_1(0 + 0.5, 4 + (-1.78125) \cdot 0.5)$
 $= -0.5 \cdot 3.109375 = -1.5546875$
 $k_{4,2} = f_2(0 + 0.5, 6 + 1.715125 \cdot 0.5)$
 $= 4 - 0.3 \cdot 6.857625 - 0.1 \cdot 3.512625 = 1.63179375$
 $\Delta y_{i+1} = y_i + \frac{1}{6} (k_1 + 2k_2 + 2k_3 + k_4)h$
 $y_1[2] = y_1[1] + \frac{1}{6} (k_1 + 2k_2 + 2k_3 + k_4)h$
 $= \begin{bmatrix} 2.115234 \\ 6.897670 \end{bmatrix}$

solving high-order ODE

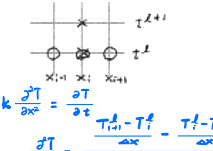
$\begin{cases} \frac{d^2 y}{dx^2} + 2 \frac{dy}{dx} = x + y \\ y(1) = \frac{dy}{dx} = 2 \text{ when } x=0 \end{cases}$
 let $u = \frac{dy}{dx}$
 $\frac{du}{dx} = \frac{d^2 y}{dx^2}$
 $\Rightarrow \frac{du}{dx} + 2u = x + y$
 $\Rightarrow \frac{du}{dx} = x + y - 2u$
 $\begin{cases} \frac{d^2 y}{dx^2} = u & y(0) = 1 \\ \frac{du}{dx} = x + y - 2u & u(0) = 2 \end{cases}$
 use this to propagate
 $\begin{cases} \frac{d^2 y}{dx^2} + 2 \frac{dy}{dx} = x + y \\ y(1) = \frac{dy}{dx} = 2 \end{cases}$
 $y(1) = 2, \frac{dy}{dx}(1) = 2$
 $\frac{d^2 y}{dx^2} = u, \frac{du}{dx} = x + y - 2u$
 $\begin{cases} \frac{dy}{dx} = u & y(0) = 1 \\ \frac{du}{dx} = v & u(0) = 2 \\ \frac{dv}{dx} = v & v(0) = 3 \\ & x=0 \end{cases}$
 use this to propagate

PDE

Laplace equation
 $\frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2} = 0$


$\frac{\partial^2 T}{\partial x^2} = \frac{(T_{i,j+1} - T_{i,j}) - (T_{i,j} - T_{i,j-1}))}{\Delta x}$
 $= \frac{T_{i,j+1} - 2T_{i,j} + T_{i,j-1}}{\Delta x^2}$
 $\frac{\partial^2 T}{\partial y^2} = \frac{(T_{i,j+1} - T_{i,j}) - (T_{i,j} - T_{i,j-1})}{\Delta y}$
 $= \frac{T_{i,j+1} - 2T_{i,j} + T_{i,j-1}}{\Delta y^2}$
 $\frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2} = 0$
 $\Rightarrow \frac{T_{i,j+1} - 2T_{i,j} + T_{i,j-1}}{\Delta x^2} + \frac{T_{i,j+1} - 2T_{i,j} + T_{i,j-1}}{\Delta y^2} = 0$
 $\Delta x = \Delta y$
 $\Rightarrow T_{i,j+1} - 2T_{i,j} + T_{i,j-1} + T_{i,j+1} - 2T_{i,j} + T_{i,j-1} - 4T_{i,j} = 0$
 $\Rightarrow T_{i,j} = \frac{T_{i,j+1} + T_{i,j-1} + T_{i,j+1} + T_{i,j-1}}{4}$

Parabolic Equation

$k \frac{\partial^2 T}{\partial x^2} = \frac{\partial T}{\partial t}$

 $k \frac{\partial^2 T}{\partial x^2} = \frac{\partial T}{\partial t}$
 $\frac{\partial T}{\partial x^2} = \frac{T_{i,j+1} - T_{i,j}^2 - T_{i,j}^2 - T_{i,j}^2}{\Delta x^2}$
 $\frac{\partial T}{\partial t} = \frac{T_{i,j+1}^n - T_{i,j}^n - T_{i,j}^n - T_{i,j}^n}{\Delta t}$
 $\Rightarrow k \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta x^2} = \frac{T_{i,j}^{n+1} - T_{i,j}^n}{\Delta t}$
 $\Rightarrow T_{i,j}^{n+1} - T_{i,j}^n = \frac{k \Delta t}{\Delta x^2} (T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n)$
 $\Rightarrow T_{i,j}^{n+1} = T_{i,j}^n + \lambda (T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n)$
 $\lambda = \frac{k \Delta t}{\Delta x^2}$
 $\Delta x = 2 \text{ cm}$
 $\Delta t = 0.1 \text{ s}$
 $k = 0.825 \text{ cm}^2/\text{s}$
 $\lambda = 0.20875$
 $\begin{cases} T(0) = 100 @ \forall t \\ T(10) = 50 \\ T(x) = 0 @ t=0, 0 < x < 10 \end{cases}$

$\frac{\partial T}{\partial t} = k \left(\frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2} \right)$

$\frac{\partial T}{\partial t} = \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta x^2} + \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta y^2}$
 $\frac{\partial T}{\partial t} = \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta x^2} + \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta y^2}$
 $\frac{\partial T}{\partial t} = \frac{T_{i,j+1}^n - T_{i,j}^n}{\Delta x} - \frac{T_{i,j}^n - T_{i,j-1}^n}{\Delta x}$
 $\frac{\partial T}{\partial t} = \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta x^2} + \frac{T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n}{\Delta y^2}$
 $\Rightarrow T_{i,j}^{n+1} = T_{i,j}^n + \frac{k \Delta t}{\Delta x^2} (T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n) + \frac{k \Delta t}{\Delta y^2} (T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n)$
 let $\Delta x = \Delta y$ $\lambda = \frac{k \Delta t}{\Delta x^2} = \frac{k \Delta t}{\Delta y^2}$
 $\Rightarrow T_{i,j}^{n+1} = T_{i,j}^n + \lambda (T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n + T_{i,j+1}^n - 2T_{i,j}^n + T_{i,j-1}^n)$

convergence

1- dimension:
 $\lambda = \frac{k \Delta t}{\Delta x^2} \leq \frac{1}{2}$
 2- dimension:
 $\Delta t \leq \frac{1}{8} \frac{\Delta x^2}{k}$
 if $\Delta x = \Delta y$
 $\Delta t \leq \frac{1}{4} \frac{\Delta x^2}{k}$

compute T @ t=0.1 s

$T_1^1 = T_0^0 + 0.20875 (T_2^0 - 2T_1^0 + T_0^0)$
 $= 0 + 0.20875 (0 - 2 \cdot 0 + 100)$
 $= 20.875$
compute T @ t=0.1 s
 $x = 4 \text{ cm}$
 $T_2^1 = T_1^1 + \lambda (T_3^1 - 2T_2^1 + T_1^1)$
 $= 0 + 0.20875 (0 - 2 \cdot 0 + 0)$
 $= 0$
compute T @ t=0.1 s
 $x = 6 \text{ cm}$
 $T_3^1 = T_2^1 + 0.20875 (T_4^1 - 2T_3^1 + T_2^1)$
 $= 0 + 0.20875 (0 - 2 \cdot 0 + 0)$
 $= 0$
compute T @ t=0.1 s
 $x = 8 \text{ cm}$
 $T_4^1 = T_3^1 + 0.20875 (T_5^1 - 2T_4^1 + T_3^1)$
 $= 0 + 0.20875 (50 - 0 + 0)$
 $= 10.4375$
compute T @ t=0.2 s
 $x = 2 \text{ cm}$
 $T_1^2 = T_1^1 + \lambda (T_2^1 - 2T_1^1 + T_0^1)$
 $= 20.875 + 0.20875 (0 - 2 \cdot 20.875 + 100)$
 $= 40.875$

convergence

$\lambda = \frac{k \Delta t}{\Delta x^2} \leq \frac{1}{4} \Rightarrow \Delta t \leq \frac{1}{4} \frac{\Delta x^2}{k}$

Linear Programming

Example problem

1. origin 1
2. origin 2

terminal 1 b_1
terminal 2 b_2
terminal 3 b_3

minimize $[C_1 \ C_2 \ C_3] \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} + [C_4 \ C_5 \ C_6] \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}$

s.t. $\begin{cases} x_1 + x_2 + x_3 = a_1 \\ x_1 + x_2 + x_3 = a_2 \\ x_1 + x_2 = b_1 \\ x_1 + x_3 = b_2 \\ x_2 + x_3 = b_3 \\ x_1, x_2, x_3 \geq 0 \end{cases}$

minimize $[C_1 \ C_2 \ C_3 \ C_4 \ C_5 \ C_6] \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ y_1 \\ y_2 \\ y_3 \end{bmatrix}$

s.t. $\begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} a_1 \\ a_2 \\ b_1 \\ b_2 \\ b_3 \\ b_3 \end{bmatrix}$

minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$

Transformation to standard form

minimize $C^T x$
s.t. $Ax \leq b$
 $x \geq 0$

minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$

minimize $C^T x$
s.t. $Ax \geq b$
 $x \geq 0$

minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$

minimize $C^T x + y$
s.t. $Ax + y = b$
 $x \geq 0$
 $(ny \geq 0)$

minimize $C^T x + u - v$
s.t. $Ax + u - v = b$
 $x \geq 0$
 $u \geq 0, v \geq 0$

minimize $C^T x + y + z$
s.t. $Ax + y + z = b$
 $y \geq 0, z \geq 0$
 $x \geq 0$

minimize $C^T x + y + b - Ax - y$
s.t. $y + b - Ax - y = d$
 $x \geq 0, y \geq 0$

minimize $C^T x + b - Ax$
s.t. $b - Ax = d$
 $x \geq 0, y \geq 0$

Solution - Simplex method
高中线性最优化

Solution - Basic Feasible solution (BFS)

minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$

let x be
1. $x \in D$, i.e., $Ax = b$
 $x \geq 0$

2. cols of A corresponding to nonzero of x are linearly independent

Theorem: x is an extreme (point) iff it is BFS

1. Basic/non-Basic variables

• for any BFS x , we rewrite:
 $x = \begin{pmatrix} x_B \\ x_N \end{pmatrix}$ B : non-zero
 $A = (B \ N)$ N : zero
 $C = \begin{pmatrix} C_B \\ C_N \end{pmatrix}$

• $Ax = b$
 $Bx_B + Nx_N = b$
 $Bx_B = b$

• $Bx_B + Nx_N = b$
 $x_B = B^{-1}(b - Nx_N) \geq 0^{m_b}$
 $x_N \geq 0^{n-m_b}$

2. minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$

minimize $C_B^T x_B + C_N^T x_N$
s.t. $Bx_B + Nx_N = b$
 $x_B \geq 0^{m_b}$
 $x_N \geq 0^{n-m_b}$

minimize $C_B^T (B^{-1}(b - Nx_N)) + C_N^T x_N$
s.t. $B^{-1}(b - Nx_N) \geq 0^{m_b}$
 $x_N \geq 0^{n-m_b}$

Simplex method - the tableau

minimize $x_1 - x_2$
s.t. $\begin{cases} x_1 - x_2 \leq 2 \\ x_1 + x_2 \leq 6 \\ x_1, x_2 \geq 0 \end{cases}$

$x_1 = 0, x_2 = 6$
 $x_3 = 0, x_4 = 8$

minimize $P = x_1 - x_2$
s.t. $\begin{cases} x_1 - x_2 + s_1 = 2 \\ x_1 + x_2 + s_2 = 6 \\ x_1, x_2, s_1, s_2 \geq 0 \end{cases}$

$-x_1 + x_2 = P$
 $-P + x_1 - x_2 = 0$
 $\Rightarrow P + x_1 - x_2 = 0$

$\begin{bmatrix} 1 & -1 & 0 & 0 & 0 & 2 \\ 1 & 1 & 0 & 0 & 0 & 6 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \\ P \end{bmatrix} = \begin{bmatrix} 2 \\ 6 \\ 0 \\ 0 \\ 0 \end{bmatrix}$

$\begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 8 \\ 1 & 1 & 0 & 0 & 0 & 6 \\ 1 & -1 & 0 & 0 & 1 & 0 \end{bmatrix}$

$\begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 8 \\ 1 & 1 & 0 & 0 & 0 & 6 \\ 0 & 0 & 0 & 1 & 1 & 6 \end{bmatrix}$

maximize $P = 5x + 4y$
s.t. $\begin{cases} 3x + 5y \leq 78 \\ 4x + y \leq 36 \\ x, y \geq 0 \end{cases}$

minimize $P = 5x + 4y$
s.t. $\begin{cases} 3x + 5y + s_1 = 78 \\ 4x + y + s_2 = 36 \\ x, y, s_1, s_2 \geq 0 \end{cases}$

$\begin{bmatrix} 3 & 5 & 1 & 0 & 0 & 78 \\ 4 & 1 & 0 & 1 & 0 & 36 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} x \\ y \\ s_1 \\ s_2 \\ P \end{bmatrix} = \begin{bmatrix} 78 \\ 36 \\ 0 \end{bmatrix}$

$\begin{bmatrix} 3 & 5 & 1 & 0 & 0 & 78 \\ 1 & 1/4 & 0 & 1/4 & 0 & 9 \\ 0 & -1/4 & 0 & 3/4 & 1 & 4.5 \end{bmatrix}$

$\begin{bmatrix} 0 & 1 & 1/4 & -3/4 & 0 & 12 \\ 1 & 1/4 & 0 & 1/4 & 0 & 9 \\ 0 & -1/4 & 0 & 3/4 & 1 & 4.5 \end{bmatrix}$

$\begin{bmatrix} 0 & 1 & 1/4 & -3/4 & 0 & 12 \\ 1 & 0 & -1/4 & 1/4 & 0 & 6 \\ 0 & 0 & 1/4 & 13/4 & 1 & 7.8 \end{bmatrix}$

$\begin{bmatrix} 0 & 1 & 1/4 & -3/4 & 0 & 12 \\ 1 & 0 & -1/4 & 1/4 & 0 & 6 \\ 0 & 0 & 1/4 & 13/4 & 1 & 7.8 \end{bmatrix}$

$x = 6$
 $y = 12$
 $P = 78$

Integer/Mixed-Integer Linear Programming

ILP/MILP

minimize $C^T x$
s.t. $Ax = b$
 $x \geq 0$
 $x \in \{x | x_i \in \mathbb{Z}, i \in M\}$

Solution - Branch & Bound

• Branch: partition D into subsets (check objective value in subsets)

• Bound: acquire a bound w.o./ integer constraints

e.g. minimize $7x_1 + 5x_2 + 6x_3$
s.t. $\begin{cases} 6x_1 + 3x_2 + 5x_3 \leq 10 \\ -x_1 + x_3 \leq 0 \\ x_i \in \{0, 1\} \quad i=1,2,3 \end{cases}$

maximize $7x_1 + 5x_2 + 6x_3$
s.t. $\begin{cases} 6x_1 + 3x_2 + 5x_3 \leq 10 \\ -x_1 + x_3 \leq 0 \\ x_i \leq 1 \quad i=1,2,3 \\ x_i \geq 0 \quad i=1,2,3 \end{cases}$

Bound:

• w.o. \geq constraints:
 $x_1 = 1$
 $x_2 = 1$
 $x_3 = 0.2$
 $P = 15.2$ (UB)

• Rounded down solution:
 $x_1 = 1$
 $x_2 = 1$
 $x_3 = 0$
 $P = 14$ (LB)

Branch

UB 15.2 (1, 1, 0.2)
LB 14 (1, 1, 0)

$x_2 = 0$ $x_3 = 1$

feasible infeasible

UB 14 (1, 1, 0)
LB 14 (1, 1, 0)

no need to check LV
UB = LB

Solutions to Optimization Problems

recall Duality

minimize $f_0(x)$
s.t. $f_i(x) \leq 0$
 $h_i(x) = 0$

Δ Lagrangian:
 $\mathcal{L}(x, \lambda, \mu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^l \mu_i h_i(x)$

Δ Lagrangian Dual function:
let $g(\lambda, \mu) = \inf_x \{ \mathcal{L}(x, \lambda, \mu) \}$

maximize $g(\lambda, \mu)$
s.t. $\lambda \geq 0$

let $\lambda^*, \mu^*, \alpha^* = g(\lambda^*, \mu^*)$
be optimal
 $P^* \geq d^*$

Δ Strong duality
Slater's condition
 $\exists x \in \text{relint } D$
s.t. $f_i(x) < 0$
 $h_i(x) = 0$

Δ complementary slackness
 $f_0(x^*) = g(\lambda^*, \mu^*)$
 $= \inf_x \{ f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^l \mu_i^* h_i(x) \}$
 $\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^l \mu_i^* h_i(x^*)$
 $\leq f_0(x^*)$

$\lambda_i^* f_i(x^*) = 0$
 $\Rightarrow \begin{cases} f_i(x^*) < 0 \rightarrow \lambda_i^* = 0 \\ \lambda_i^* > 0 \rightarrow f_i(x^*) = 0 \end{cases}$

Δ KKT optimality conditions

$\begin{cases} f_i(x) \leq 0 \\ h_i(x) = 0 \\ \lambda_i^* \geq 0 \\ \lambda_i^* f_i(x^*) = 0 \end{cases}$

$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^l \mu_i^* \nabla h_i(x^*) = 0$

this will be used later

Descent Method

1. given $x^{(0)} \in \text{dom } f$

2. repeat
1) descent direction Δx
2) stepsize $t \geq 0$
3) $x^{k+1} = x^k + t^k \Delta x^k$

3. Stop
 $\| \nabla f(x) \|_2 \leq \epsilon$

Line Search

Exact Line Search
 $t = \text{argmin}_{t \geq 0} f(x + t \Delta x)$

Backtracking Line Search
1. given Δx^k
 $f(x^k)$
 $\alpha \in (0, 0.5)$
 $\beta \in (0, 1)$

2. set $t = 1$

3. while $f(x + t \Delta x) > f(x) + \alpha t \nabla f(x)^T \Delta x$
 $t = \beta t$

Primal-dual
Interior-point method
minimize $f_0(x) + \beta(x)$
s.t. $Ax = b$

KKT:
 $\begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \frac{1}{\beta(x^*(t))} \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax = b \\ \nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} \nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + A^T \bar{w} = 0 \\ -\lambda_i \nabla f_i(x) = \lambda_i \nabla f_i(x) \\ Ax = b \end{cases}$

$\Rightarrow \lambda_c(x, \lambda, \mu) = \begin{bmatrix} \nabla f_0(x) + Df(x)^T \lambda + A^T \mu \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix}$

Newton's method
 $\Rightarrow \lambda_c(x + \Delta x, \lambda + \Delta \lambda, \mu + \Delta \mu) = \begin{bmatrix} \nabla f_0(x) + Df(x)^T \lambda + A^T \mu \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix} + \begin{bmatrix} \nabla^2 f_0(x) & D^2 f_0(x)^T \lambda & D^2 f_0(x)^T \mu \\ \nabla^2 f_i(x) & \nabla^2 f_i(x) & \nabla^2 f_i(x)^T \mu \\ A & 0 & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta \lambda \\ \Delta \mu \end{bmatrix} + \begin{bmatrix} \nabla f_0(x) \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix}$

Unconstrained problem

Gradient descent

1. given $x_0 \in \text{dom } f$

2. repeat
 $\Delta x^k = -\nabla f(x^k)$
 t : from line search
 $x^{k+1} = x^k + t \Delta x^k$

3. until $\| \nabla f(x) \| \leq \epsilon$

Newton's method

1. given $x^0 \in \text{dom } f$

2. repeat:
 $\Delta x^k = -\nabla^2 f(x^k)^{-1} \nabla f(x^k)$
 t : from line search
 $x^{k+1} = x^k + t \Delta x^k$

3. until $\lambda_2^k \leq \epsilon$

on Newton's Method
 $f(x)$
 $\Rightarrow f'(x+v) = f'(x) + \nabla^2 f(x)^T v + \frac{1}{2} \nabla^3 f(x)^T v^2$
 $\Rightarrow g f'(x+v) = 0$
 $\Rightarrow \nabla^2 f(x) + \nabla^3 f(x)^T v = 0$
 $\Rightarrow v = -\nabla^2 f(x)^{-1} \nabla f(x)$

Equality-constrained Problem

minimize $f(x)$
s.t. $Ax = b$

relax
Slater's condition
 $\exists x \in \text{relint } D$
s.t. $Ax = b$
KKT optimality condition

$f(x) \leq 0$
 $h_i(x) = 0$
 $\lambda_i^* \geq 0$
 $\lambda_i^* f_i(x^*) = 0$

$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^l \mu_i^* \nabla h_i(x^*) = 0$

Δ Here @ optimal x, \bar{w}
 $A \bar{x} = b$
 $\nabla f_0(\bar{x}) + A^T \bar{w} = 0$

Newton's Method

1. @ x
addition to x : $x + v$
 $\therefore x + v \in \text{dom } f$
 $A(x+v) = b$
 $\therefore Av = 0$ $\text{---} \textcircled{1}$

2. $f(x+v)$
 $= f(x) + \nabla f(x)^T v + \frac{1}{2} \nabla^2 f(x)^T v^2$
s.t. $Av = 0$ $\text{---} \textcircled{2}$

Δ \bar{v} is optimal of $\textcircled{2}$
 $\therefore \exists \bar{w}$ s.t.
 $\nabla f(x) + \nabla^2 f(x)^T \bar{v} + A^T \bar{w} = 0$
and recall $A \bar{v} = 0$
from KKT $\mathcal{L}(v, w)$
 $= f(x) + \nabla f(x)^T v + \frac{1}{2} \nabla^2 f(x)^T v^2 + \bar{w}^T Av$
 $\therefore A \bar{v} = 0$
 $\nabla f(x) + \nabla^2 f(x)^T \bar{v} + A^T \bar{w} = 0$
 $\Rightarrow \begin{bmatrix} \nabla f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \bar{v} \\ \bar{w} \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$
 $\therefore \bar{v} = \Delta x_{\text{nt}}$

Inequality-Constrained Problem

Barrier Method

minimize $f_0(x)$
s.t. $f_i(x) \leq 0$
 $Ax = b$

minimize $t f_0(x) + \beta(x)$
s.t. $Ax = b$
 $\beta(x) = -\sum_{i=1}^m \lambda_i g_i(-f_i(x))$

1. given $x^0 \in \text{dom } f$
 $t^0 := t^0 > 0$
 $\mu > 1$
 $\epsilon > 0$

2. repeat
a. solve
minimize $t f_0(x) + \beta(x)$
s.t. $Ax = b$
(get x^k, w^k, t^k)
b. $x = x^k(t)$
c. until $\| \lambda_i^k \| < \epsilon$
d. $t := \mu t$

Interior-point method

minimize $f_0(x)$
s.t. $f_i(x) \leq 0$
 $Ax = b$

KKT:
 $\begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \frac{1}{\beta(x^*(t))} \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax = b \\ \nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \lambda_c(x, \lambda, \mu) = \begin{bmatrix} \nabla f_0(x) + Df(x)^T \lambda + A^T \mu \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix}$

Newton's method
 $\Rightarrow \lambda_c(x + \Delta x, \lambda + \Delta \lambda, \mu + \Delta \mu) = \begin{bmatrix} \nabla f_0(x) + Df(x)^T \lambda + A^T \mu \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix} + \begin{bmatrix} \nabla^2 f_0(x) & D^2 f_0(x)^T \lambda & D^2 f_0(x)^T \mu \\ \nabla^2 f_i(x) & \nabla^2 f_i(x) & \nabla^2 f_i(x)^T \mu \\ A & 0 & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta \lambda \\ \Delta \mu \end{bmatrix} + \begin{bmatrix} \nabla f_0(x) \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix}$

Primal-dual
Interior-point method
minimize $f_0(x) + \beta(x)$
s.t. $Ax = b$

KKT:
 $\begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \frac{1}{\beta(x^*(t))} \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax = b \\ \nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \lambda_c(x, \lambda, \mu) = \begin{bmatrix} \nabla f_0(x) + Df(x)^T \lambda + A^T \mu \\ -\lambda_i \nabla f_i(x) \\ Ax - b \end{bmatrix}$

Newton's method
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Barrier Method

minimize $f_0(x)$
s.t. $f_i(x) \leq 0$
 $Ax = b$

minimize $t f_0(x) + \beta(x)$
s.t. $Ax = b$
 $\beta(x) = -\sum_{i=1}^m \lambda_i g_i(-f_i(x))$

1. given $x^0 \in \text{dom } f$
 $t^0 := t^0 > 0$
 $\mu > 1$
 $\epsilon > 0$

2. repeat
a. solve
minimize $t f_0(x) + \beta(x)$
s.t. $Ax = b$
(get x^k, w^k, t^k)
b. $x = x^k(t)$
c. until $\| \lambda_i^k \| < \epsilon$
d. $t := \mu t$

Primal-dual
Interior-point method
minimize $f_0(x) + \beta(x)$
s.t. $Ax = b$

KKT:
 $\begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

$\Rightarrow \begin{cases} Ax^*(t) = b \\ \nabla f_0(x^*(t)) + \frac{1}{\beta(x^*(t))} \nabla \beta(x^*(t)) + A^T \bar{w} = 0 \end{cases}$

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